



COMP 512  
Rice University  
Spring 2013

*THE PARTITIONING ALGORITHM FOR DETECTING  
CONGRUENT EXPRESSIONS*

*YET ANOTHER WAY TO ACHIEVE REDUNDANCY ELIMINATION*

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## BACKGROUND

This algorithm discovers *congruent expressions*, which should have the same value at runtime. Sets of congruent expressions can serve as the basis for a powerful form of redundancy elimination.

Critical points:

- The Partitioning Method
  - ◆ Using a well-known algorithm for a new purpose
  - ◆ Classic example of an optimistic algorithm
  - ◆ Performs analysis not transformation
- Compiler must still rewrite the code
  - ◆ Paper suggests a method based on dominance
  - ◆ Better options are available

We want to cover this technique, in part, as a prelude to the lecture on algebraic reassociation of expressions.



## PARTITIONING METHOD

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Key idea

Operations  $x$  and  $y$  are congruent iff they have the same operator & their operands are congruent

Congruence is related to redundancy, but not the same

- ◆ *Redundancy implicitly includes a notion that  $x$  and  $y$  share an execution path*
- ◆ *Congruence defers that issue for later consideration*

Using the idea

- Start with SSA form (we need the name space)
- Partition the set of all expressions into congruence classes
  - ◆ Find largest possible sets
- Each class needs one (*or more*) representative computation
- Other uses can be replaced with the representer



## PARTITIONING METHOD

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### The algorithm

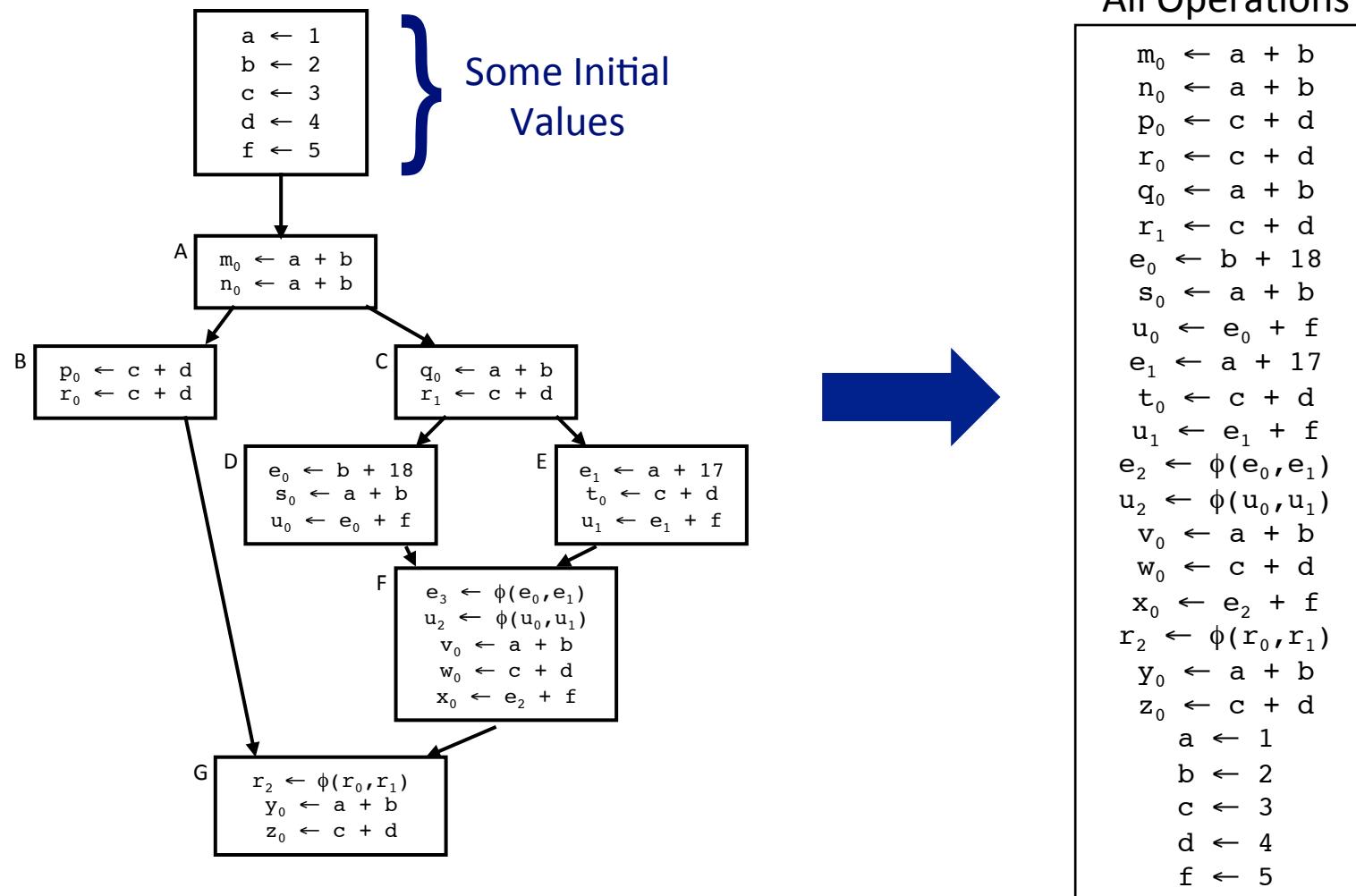
1. Partition operations into initial classes by opcode  
*Treat constant values as unique opcodes*
2.  $Worklist \leftarrow$  all classes
3. While  $Worklist$  is not empty
  - A. Remove a class  $c$  from  $Worklist$
  - B. For each class  $s$  that uses some  $x$  defined in  $c$ 
    - i. Split  $s$  into  $s_1$  and  $s_2$  around uses of  $c$
    - ii. Remove  $s$  from  $Worklist$  & add the smaller of  $s_1$  and  $s_2$  to  $Worklist$
4. Pick a representer for each class & rewrite the code to replace other uses of the class with representer

Based on Hopcroft's algorithm for DFA minimization

- ◆ *A widely-used, often re-invented algorithm*



## PARTITIONING METHOD





# PARTITIONING METHOD

```
m0 ← a + b  
n0 ← a + b  
p0 ← c + d  
r0 ← c + d  
q0 ← a + b  
r1 ← c + d  
e0 ← b + 18  
s0 ← a + b  
u0 ← e0 + f  
e1 ← a + 17  
t0 ← c + d  
u1 ← e1 + f  
e2 ← φ(e0, e1)  
u2 ← φ(u0, u1)  
v0 ← a + b  
w0 ← c + d  
x0 ← e2 + f  
r2 ← φ(r0, r1)  
y0 ← a + b  
z0 ← c + d  
a ← 1  
b ← 2  
c ← 3  
d ← 4  
f ← 5
```

Compute  
Initial Partition

1	a ← 1
2	b ← 2
3	c ← 3
4	d ← 4
5	f ← 5
6	e <sub>2</sub> ← φ(e <sub>0</sub> , e <sub>1</sub> ) u <sub>2</sub> ← φ(u <sub>0</sub> , u <sub>1</sub> ) r <sub>2</sub> ← φ(r <sub>0</sub> , r <sub>1</sub> )

7	m <sub>0</sub> ← a + b n <sub>0</sub> ← a + b p <sub>0</sub> ← c + d r <sub>0</sub> ← c + d q <sub>0</sub> ← a + b r <sub>1</sub> ← c + d e <sub>0</sub> ← b + 18 s <sub>0</sub> ← a + b u <sub>0</sub> ← e <sub>0</sub> + f e <sub>1</sub> ← a + 17 t <sub>0</sub> ← c + d u <sub>1</sub> ← e <sub>1</sub> + f v <sub>0</sub> ← a + b w <sub>0</sub> ← c + d x <sub>0</sub> ← e <sub>2</sub> + f Y <sub>0</sub> ← a + b z <sub>0</sub> ← c + d
---	---

Worklist  
1, 2, 3, 4, 5, 6, 7

# PARTITIONING METHOD



```

1   a ← 1
2   b ← 2
3   c ← 3
4   d ← 4
5   f ← 5
6   e2 ← φ(e0, e1)
    u2 ← φ(u0, u1)
    r2 ← φ(r0, r1)

```

7

```

m0 ← a + b
n0 ← a + b
p0 ← c + d
r0 ← c + d
q0 ← a + b
r1 ← c + d
e0 ← b + 18
s0 ← a + b
u0 ← e0 + f
e1 ← a + 17
t0 ← c + d
u1 ← e1 + f
v0 ← a + b
w0 ← c + d
x0 ← e2 + f
y0 ← a + b
z0 ← c + d

```

Split  
on a

```

1   a ← 1
2   b ← 2
3   c ← 3
4   d ← 4
5   f ← 5
6   e2 ← φ(e0, e1)
    u2 ← φ(u0, u1)
    r2 ← φ(r0, r1)

```

7

```

m0 ← a + b
n0 ← a + b
q0 ← a + b
s0 ← a + b
e1 ← a + 17
v0 ← a + b
y0 ← a + b

```

8

```

p0 ← c + d
r0 ← c + d
r1 ← c + d
e0 ← b + 18
u0 ← e0 + f
t0 ← c + d
u1 ← e1 + f
w0 ← c + d
x0 ← e2 + f
z0 ← c + d

```

Worklist:  
2, 3, 4, 5, 6, 7



## PARTITIONING METHOD

1	$a \leftarrow 1$
2	$b \leftarrow 2$
3	$c \leftarrow 3$
4	$d \leftarrow 4$
5	$f \leftarrow 5$
6	$e_2 \leftarrow \phi(e_0, e_1)$ $u_2 \leftarrow \phi(u_0, u_1)$ $r_2 \leftarrow \phi(r_0, r_1)$

7	$m_0 \leftarrow a + b$ $n_0 \leftarrow a + b$ $q_0 \leftarrow a + b$ $s_0 \leftarrow a + b$ $e_1 \leftarrow a + 17$ $v_0 \leftarrow a + b$ $y_0 \leftarrow a + b$
8	$p_0 \leftarrow c + d$ $r_0 \leftarrow c + d$ $r_1 \leftarrow c + d$ $e_0 \leftarrow b + 18$ $u_0 \leftarrow e_0 + f$ $t_0 \leftarrow c + d$ $u_1 \leftarrow e_1 + f$ $w_0 \leftarrow c + d$ $x_0 \leftarrow e_2 + f$ $z_0 \leftarrow c + d$

Split  
on b

1	$a \leftarrow 1$
2	$b \leftarrow 2$
3	$c \leftarrow 3$
4	$d \leftarrow 4$
5	$f \leftarrow 5$
6	$e_2 \leftarrow \phi(e_0, e_1)$ $u_2 \leftarrow \phi(u_0, u_1)$ $r_2 \leftarrow \phi(r_0, r_1)$

7	$m_0 \leftarrow a + b$ $n_0 \leftarrow a + b$ $q_0 \leftarrow a + b$ $s_0 \leftarrow a + b$ $v_0 \leftarrow a + b$ $y_0 \leftarrow a + b$
9	$e_1 \leftarrow a + 17$
8	$p_0 \leftarrow c + d$ $r_0 \leftarrow c + d$ $r_1 \leftarrow c + d$ $u_0 \leftarrow e_0 + f$ $t_0 \leftarrow c + d$ $u_1 \leftarrow e_1 + f$ $w_0 \leftarrow c + d$ $x_0 \leftarrow e_2 + f$ $z_0 \leftarrow c + d$
10	$e_0 \leftarrow b + 18$

Worklist  
3, 4, 5, 6, 9, 10



## PARTITIONING METHOD

```

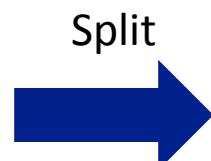
1   a ← 1
2   b ← 2
3   c ← 3
4   d ← 4
5   f ← 5
6   e2 ← φ(e0, e1)
    u2 ← φ(u0, u1)
    r2 ← φ(r0, r1)

```

```

7   m0 ← a + b
    n0 ← a + b
    q0 ← a + b
    s0 ← a + b
    v0 ← a + b
    y0 ← a + b
9   e1 ← a + 17
8   p0 ← c + d
    r0 ← c + d
    r1 ← c + d
    u0 ← e0 + f
    t0 ← c + d
    u1 ← e1 + f
    w0 ← c + d
    x0 ← e2 + f
    z0 ← c + d
10  e0 ← b + 18

```



on c

Worklist  
4, 5, 6, 9, 10, 11

```

1   a ← 1
2   b ← 2
3   c ← 3
4   d ← 4
5   f ← 5
9   e1 ← a + 17
6   e2 ← φ(e0, e1)
    u2 ← φ(u0, u1)
    r2 ← φ(r0, r1)

```

```

7   m0 ← a + b
    n0 ← a + b
    q0 ← a + b
    s0 ← a + b
    v0 ← a + b
    y0 ← a + b
8   p0 ← c + d
    r0 ← c + d
    r1 ← c + d
    t0 ← c + d
    w0 ← c + d
    z0 ← c + d
11  u0 ← e0 + f
    u1 ← e1 + f
    x0 ← e2 + f
10  e0 ← b + 18

```

# PARTITIONING METHOD



1	$a \leftarrow 1$
2	$b \leftarrow 2$
3	$c \leftarrow 3$
4	$d \leftarrow 4$
5	$f \leftarrow 5$
6	$e_2 \leftarrow \phi(e_0, e_1)$ $u_2 \leftarrow \phi(u_0, u_1)$ $r_2 \leftarrow \phi(r_0, r_1)$

7	$m_0 \leftarrow a + b$ $n_0 \leftarrow a + b$ $q_0 \leftarrow a + b$ $s_0 \leftarrow a + b$ $v_0 \leftarrow a + b$ $y_0 \leftarrow a + b$
9	$e_1 \leftarrow a + 17$
8	$p_0 \leftarrow c + d$ $r_0 \leftarrow c + d$ $r_1 \leftarrow c + d$ $t_0 \leftarrow c + d$ $w_0 \leftarrow c + d$ $z_0 \leftarrow c + d$
11	$u_0 \leftarrow e_0 + f$ $u_1 \leftarrow e_1 + f$ $x_0 \leftarrow e_2 + f$
10	$e_0 \leftarrow b + 18$

Split  
on d, f

Neither d nor f refine  
the partitions

Worklist  
6, 9, 10, 11

1	$a \leftarrow 1$
2	$b \leftarrow 2$
3	$c \leftarrow 3$
4	$d \leftarrow 4$
5	$f \leftarrow 5$
6	$e_2 \leftarrow \phi(e_0, e_1)$ $u_2 \leftarrow \phi(u_0, u_1)$ $r_2 \leftarrow \phi(r_0, r_1)$

7	$m_0 \leftarrow a + b$ $n_0 \leftarrow a + b$ $q_0 \leftarrow a + b$ $s_0 \leftarrow a + b$ $v_0 \leftarrow a + b$ $y_0 \leftarrow a + b$
9	$e_1 \leftarrow a + 17$
8	$p_0 \leftarrow c + d$ $r_0 \leftarrow c + d$ $r_1 \leftarrow c + d$ $t_0 \leftarrow c + d$ $w_0 \leftarrow c + d$ $z_0 \leftarrow c + d$
11	$u_0 \leftarrow e_0 + f$ $u_1 \leftarrow e_1 + f$ $x_0 \leftarrow e_2 + f$
10	$e_0 \leftarrow b + 18$



## PARTITIONING METHOD

1     $a \leftarrow 1$   
2     $b \leftarrow 2$   
3     $c \leftarrow 3$   
4     $d \leftarrow 4$   
5     $f \leftarrow 5$   
6     $e_2 \leftarrow \phi(e_0, e_1)$   
       $u_2 \leftarrow \phi(u_0, u_1)$   
       $r_2 \leftarrow \phi(r_0, r_1)$

7     $m_0 \leftarrow a + b$   
       $n_0 \leftarrow a + b$   
       $q_0 \leftarrow a + b$   
       $s_0 \leftarrow a + b$   
       $v_0 \leftarrow a + b$   
       $y_0 \leftarrow a + b$   
  
9     $e_1 \leftarrow a + 17$   
  
8     $p_0 \leftarrow c + d$   
       $r_0 \leftarrow c + d$   
       $r_1 \leftarrow c + d$   
       $t_0 \leftarrow c + d$   
       $w_0 \leftarrow c + d$   
       $z_0 \leftarrow c + d$   
  
11    $u_0 \leftarrow e_0 + f$   
       $u_1 \leftarrow e_1 + f$   
       $x_0 \leftarrow e_2 + f$   
  
10    $e_0 \leftarrow b + 18$

Split  
on  $e_2, u_2, r_2$

Worklist  
9, 10, 11

1     $a \leftarrow 1$   
2     $b \leftarrow 2$   
3     $c \leftarrow 3$   
4     $d \leftarrow 4$   
5     $f \leftarrow 5$   
6     $e_2 \leftarrow \phi(e_0, e_1)$   
       $u_2 \leftarrow \phi(u_0, u_1)$   
       $r_2 \leftarrow \phi(r_0, r_1)$

7     $m_0 \leftarrow a + b$   
       $n_0 \leftarrow a + b$   
       $q_0 \leftarrow a + b$   
       $s_0 \leftarrow a + b$   
       $v_0 \leftarrow a + b$   
       $y_0 \leftarrow a + b$   
  
9     $e_1 \leftarrow a + 17$   
  
8     $p_0 \leftarrow c + d$   
       $r_0 \leftarrow c + d$   
       $r_1 \leftarrow c + d$   
       $t_0 \leftarrow c + d$   
       $w_0 \leftarrow c + d$   
       $z_0 \leftarrow c + d$   
  
11    $x_0 \leftarrow e_2 + f$   
  
12    $u_0 \leftarrow e_0 + f$   
       $u_1 \leftarrow e_1 + f$   
  
10    $e_0 \leftarrow b + 18$



## PARTITIONING METHOD

```

1   a ← 1
2   b ← 2
3   c ← 3
4   d ← 4
5   f ← 5
6   e2 ← φ(e0, e1)
    u2 ← φ(u0, u1)
    r2 ← φ(r0, r1)

```

```

7   m0 ← a + b
    n0 ← a + b
    q0 ← a + b
    s0 ← a + b
    v0 ← a + b
    y0 ← a + b
9   e1 ← a + 17
8   p0 ← c + d
    r0 ← c + d
    r1 ← c + d
    t0 ← c + d
    w0 ← c + d
    z0 ← c + d
11  x0 ← e2 + f
12  u0 ← e0 + f
    u1 ← e1 + f
10  e0 ← b + 18

```

Split  
on e<sub>1</sub>

Worklist  
10, 11, 12, 6

```

1   a ← 1
2   b ← 2
3   c ← 3
4   d ← 4
5   f ← 5
6   e2 ← φ(e0, e1)
    u2 ← φ(u0, u1)
    r2 ← φ(r0, r1)

```

```

7   m0 ← a + b
    n0 ← a + b
    q0 ← a + b
    s0 ← a + b
    v0 ← a + b
    y0 ← a + b
9   e1 ← a + 17
8   p0 ← c + d
    r0 ← c + d
    r1 ← c + d
    t0 ← c + d
    w0 ← c + d
    z0 ← c + d
11  x0 ← e2 + f
12  u1 ← e1 + f
15  u0 ← e0 + f
10  e0 ← b + 18

```



## PARTITIONING METHOD

```

1   a ← 1
2   b ← 2
3   c ← 3
4   d ← 4
5   f ← 5
6   e2 ← φ(e0, e1)
14 u2 ← φ(u0, u1)
r2 ← φ(r0, r1)

```

```

7   m0 ← a + b
n0 ← a + b
q0 ← a + b
s0 ← a + b
v0 ← a + b
y0 ← a + b
9   e1 ← a + 17
8   p0 ← c + d
r0 ← c + d
r1 ← c + d
t0 ← c + d
w0 ← c + d
z0 ← c + d
11  x0 ← e2 + f
12  u1 ← e1 + f
15  u0 ← e0 + f
10  e0 ← b + 18

```

Split  
on e<sub>0</sub>

Does not refine the partitions ...

Worklist  
11, 12, 6

```

1   a ← 1
2   b ← 2
3   c ← 3
4   d ← 4
5   f ← 5
6   e2 ← φ(e0, e1)
14 u2 ← φ(u0, u1)
r2 ← φ(r0, r1)

```

```

7   m0 ← a + b
n0 ← a + b
q0 ← a + b
s0 ← a + b
v0 ← a + b
y0 ← a + b
9   e1 ← a + 17
8   p0 ← c + d
r0 ← c + d
r1 ← c + d
t0 ← c + d
w0 ← c + d
z0 ← c + d
11  x0 ← e2 + f
12  u1 ← e1 + f
14  u0 ← e0 + f
10  e0 ← b + 18

```



## PARTITIONING METHOD

```

1   a ← 1
2   b ← 2
3   c ← 3
4   d ← 4
5   f ← 5
6   e2 ← φ(e0, e1)
14 u2 ← φ(u0, u1)
r2 ← φ(r0, r1)

```

```

7   m0 ← a + b
n0 ← a + b
q0 ← a + b
s0 ← a + b
v0 ← a + b
y0 ← a + b
9   e1 ← a + 17
8   p0 ← c + d
r0 ← c + d
r1 ← c + d
t0 ← c + d
w0 ← c + d
z0 ← c + d
11  x0 ← e2 + f
12  u1 ← e1 + f
15  u0 ← e0 + f
10  e0 ← b + 18

```

Split  
on  $x_0$

Does not refine the partitions ...

Worklist  
12, 6

```

1   a ← 1
2   b ← 2
3   c ← 3
4   d ← 4
5   f ← 5
6   e2 ← φ(e0, e1)
14 u2 ← φ(u0, u1)
r2 ← φ(r0, r1)

```

```

7   m0 ← a + b
n0 ← a + b
q0 ← a + b
s0 ← a + b
v0 ← a + b
y0 ← a + b
9   e1 ← a + 17
8   p0 ← c + d
r0 ← c + d
r1 ← c + d
t0 ← c + d
w0 ← c + d
z0 ← c + d
11  x0 ← e2 + f
12  u1 ← e1 + f
15  u0 ← e0 + f
10  e0 ← b + 18

```



## PARTITIONING METHOD

```

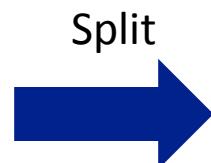
1   a ← 1
2   b ← 2
3   c ← 3
4   d ← 4
5   f ← 5
6   e2 ← φ(e0, e1)
14  u2 ← φ(u0, u1)
      r2 ← φ(r0, r1)

```

```

7   m0 ← a + b
    n0 ← a + b
    q0 ← a + b
    s0 ← a + b
    v0 ← a + b
    y0 ← a + b
9   e1 ← a + 17
8   p0 ← c + d
    r0 ← c + d
    r1 ← c + d
    t0 ← c + d
    w0 ← c + d
    z0 ← c + d
11  x0 ← e2 + f
12  u1 ← e1 + f
15  u0 ← e0 + f
10  e0 ← b + 18

```



on  $u_1$

Worklist  
6, 14

```

1   a ← 1
2   b ← 2
3   c ← 3
4   d ← 4
5   f ← 5
6   e2 ← φ(e0, e1)
14  u2 ← φ(u0, u1)
      r2 ← φ(r0, r1)

```

```

7   m0 ← a + b
    n0 ← a + b
    q0 ← a + b
    s0 ← a + b
    v0 ← a + b
    y0 ← a + b
9   e1 ← a + 17
8   p0 ← c + d
    r0 ← c + d
    r1 ← c + d
    t0 ← c + d
    w0 ← c + d
    z0 ← c + d
11  x0 ← e2 + f
12  u1 ← e1 + f
15  u0 ← e0 + f
10  e0 ← b + 18

```



## PARTITIONING METHOD

```
1   a ← 1
2   b ← 2
3   c ← 3
4   d ← 4
5   f ← 5
6   e2 ← φ(e0, e1)
14  u2 ← φ(u0, u1)
16  r2 ← φ(r0, r1)
```

```
7   m0 ← a + b
n0 ← a + b
q0 ← a + b
s0 ← a + b
v0 ← a + b
y0 ← a + b
9   e1 ← a + 17
8   p0 ← c + d
r0 ← c + d
r1 ← c + d
t0 ← c + d
w0 ← c + d
z0 ← c + d
11  x0 ← e2 + f
12  u1 ← e1 + f
15  u0 ← e0 + f
10  e0 ← b + 18
```

Split  
on e<sub>2</sub>

Does not refine the partitions ...

Worklist  
14

```
1   a ← 1
2   b ← 2
3   c ← 3
4   d ← 4
5   f ← 5
6   e2 ← φ(e0, e1)
14  u2 ← φ(u0, u1)
16  r2 ← φ(r0, r1)
```

```
7   m0 ← a + b
n0 ← a + b
q0 ← a + b
s0 ← a + b
v0 ← a + b
y0 ← a + b
9   e1 ← a + 17
8   p0 ← c + d
r0 ← c + d
r1 ← c + d
t0 ← c + d
w0 ← c + d
z0 ← c + d
11  x0 ← e2 + f
12  u1 ← e1 + f
15  u0 ← e0 + f
10  e0 ← b + 18
```



## PARTITIONING METHOD

```
1   a ← 1
2   b ← 2
3   c ← 3
4   d ← 4
5   f ← 5
6   e2 ← φ(e0, e1)
14  u2 ← φ(u0, u1)
16  r2 ← φ(r0, r1)
```

```
7   m0 ← a + b
n0 ← a + b
q0 ← a + b
s0 ← a + b
v0 ← a + b
y0 ← a + b
9   e1 ← a + 17
8   p0 ← c + d
r0 ← c + d
r1 ← c + d
t0 ← c + d
w0 ← c + d
z0 ← c + d
11  x0 ← e2 + f
12  u1 ← e1 + f
15  u0 ← e0 + f
10  e0 ← b + 18
```

Split  
on u<sub>2</sub>

Does not refine the partitions ...

Worklist is empty

```
1   a ← 1
2   b ← 2
3   c ← 3
4   d ← 4
5   f ← 5
9   e1 ← a + 17
6   e2 ← φ(e0, e1)
14  u2 ← φ(u0, u1)
16  r2 ← φ(r0, r1)
```

```
7   m0 ← a + b
n0 ← a + b
q0 ← a + b
s0 ← a + b
v0 ← a + b
y0 ← a + b
8   p0 ← c + d
r0 ← c + d
r1 ← c + d
t0 ← c + d
w0 ← c + d
z0 ← c + d
11  x0 ← e2 + f
12  u1 ← e1 + f
15  u0 ← e0 + f
10  e0 ← b + 18
```



## PARTITIONING METHOD

Now what?

1     $a \leftarrow 1$

2     $b \leftarrow 2$

3     $c \leftarrow 3$

4     $d \leftarrow 4$

5     $f \leftarrow 5$

6     $e_2 \leftarrow \phi(e_0, e_1)$

14     $u_2 \leftarrow \phi(u_0, u_1)$

16     $r_2 \leftarrow \phi(r_0, r_1)$

7     $m_0 \leftarrow a + b$   
 $n_0 \leftarrow a + b$   
 $q_0 \leftarrow a + b$   
 $s_0 \leftarrow a + b$   
 $v_0 \leftarrow a + b$   
 $y_0 \leftarrow a + b$

9     $e_1 \leftarrow a + 17$

8     $p_0 \leftarrow c + d$   
 $r_0 \leftarrow c + d$   
 $r_1 \leftarrow c + d$   
 $t_0 \leftarrow c + d$   
 $w_0 \leftarrow c + d$   
 $z_0 \leftarrow c + d$

11     $x_0 \leftarrow e_2 + f$

12     $u_1 \leftarrow e_1 + f$

15     $u_0 \leftarrow e_0 + f$

10     $e_0 \leftarrow b + 18$

We have proved a set of theorems

- Must rewrite code to use them
- Knowledge alone won't make the code run faster!

Use the partition to rebuild the code

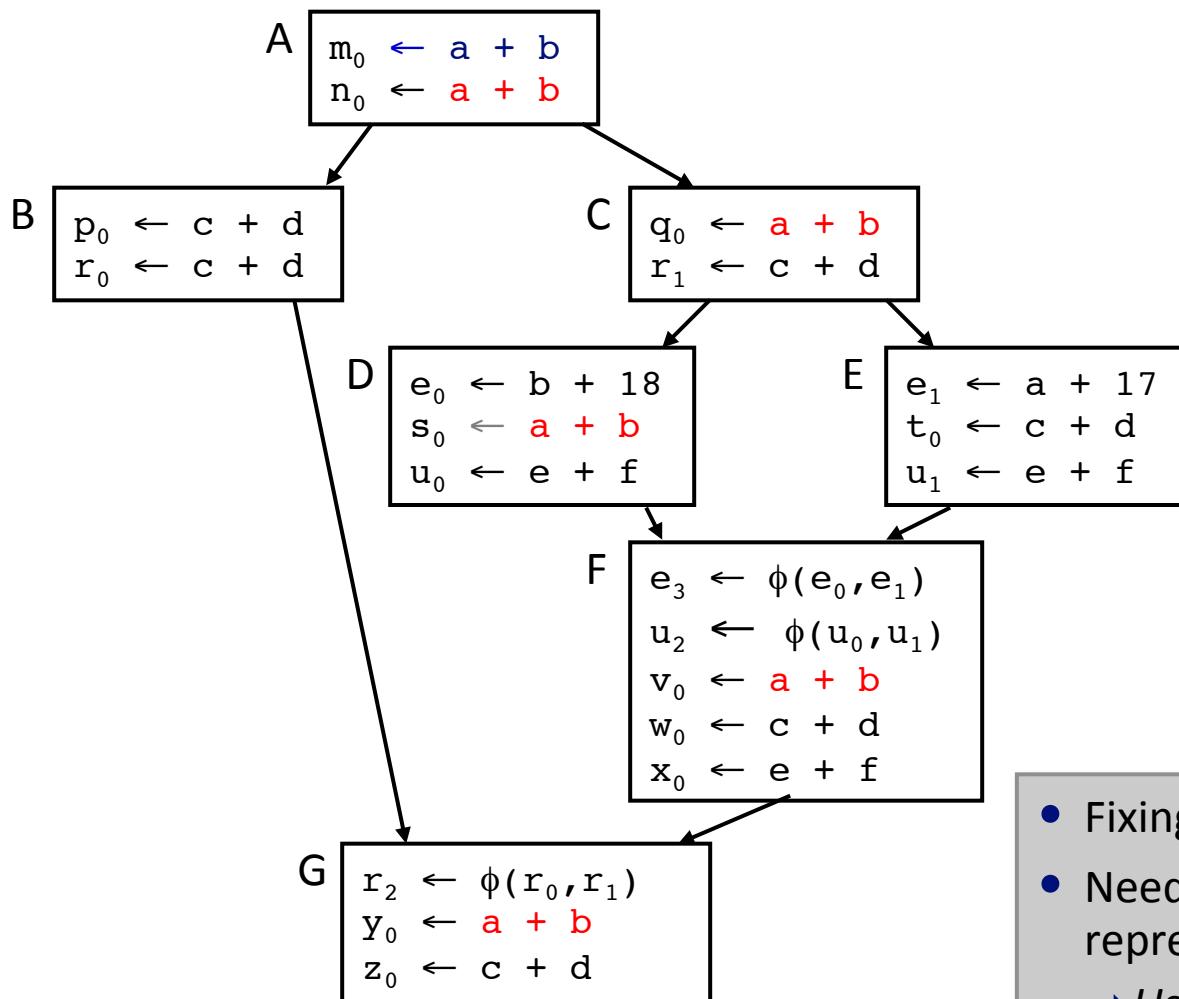
- Sets with > 1 member  $\Rightarrow$  reuse
- Sets with 1 member  $\Rightarrow$  no reuse

Let's look at the code

- Opportunities are  $a+b$  &  $c+d$
- Can we make this work?

## PARTITIONING METHOD

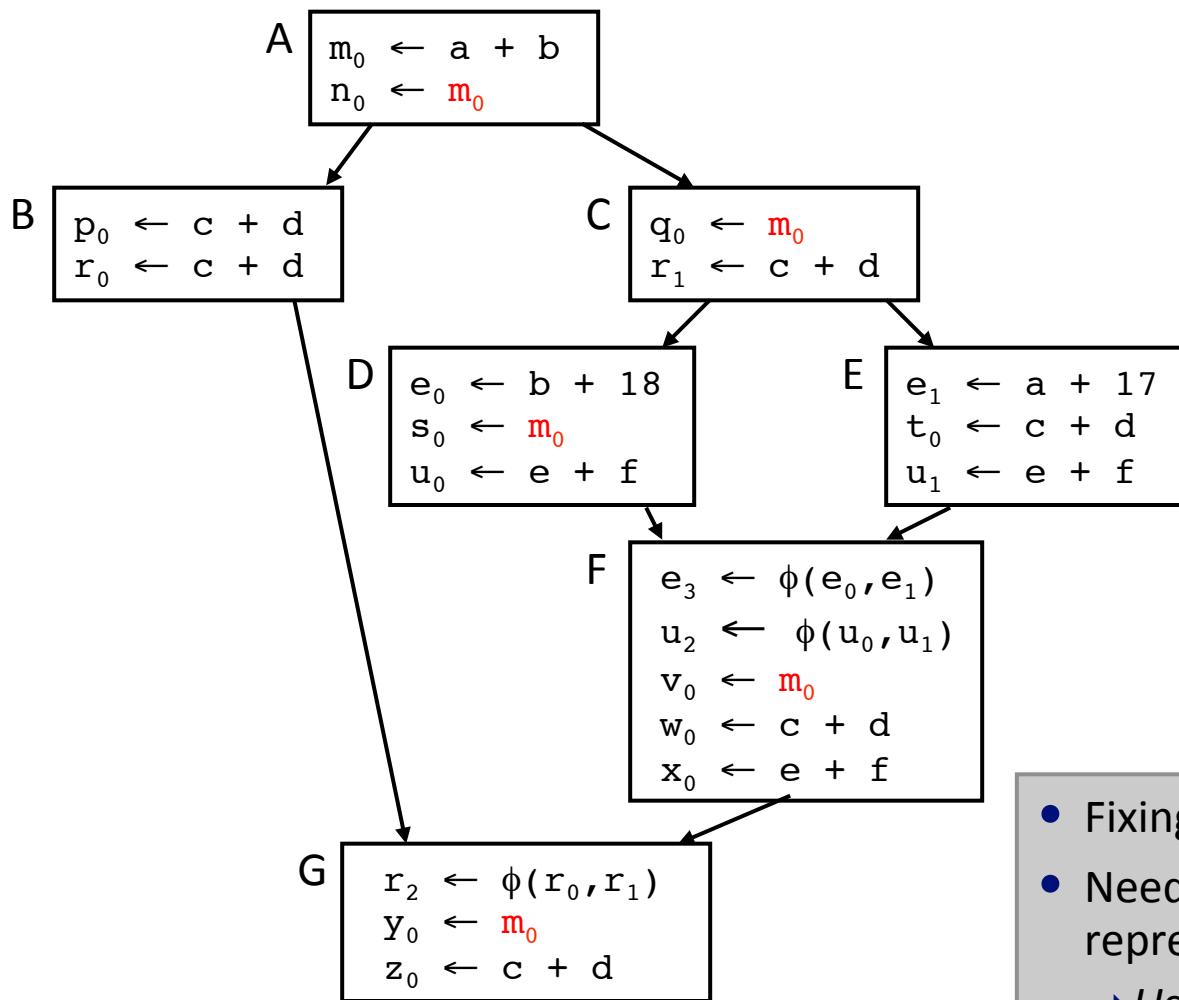
These slides omits the initialization block so that it can fit into the available space.



- Fixing  $a+b$  is easy
- Need to choose the right representer
  - Use the instance of  $a+b$  in A
  - It makes the others redundant

## PARTITIONING METHOD

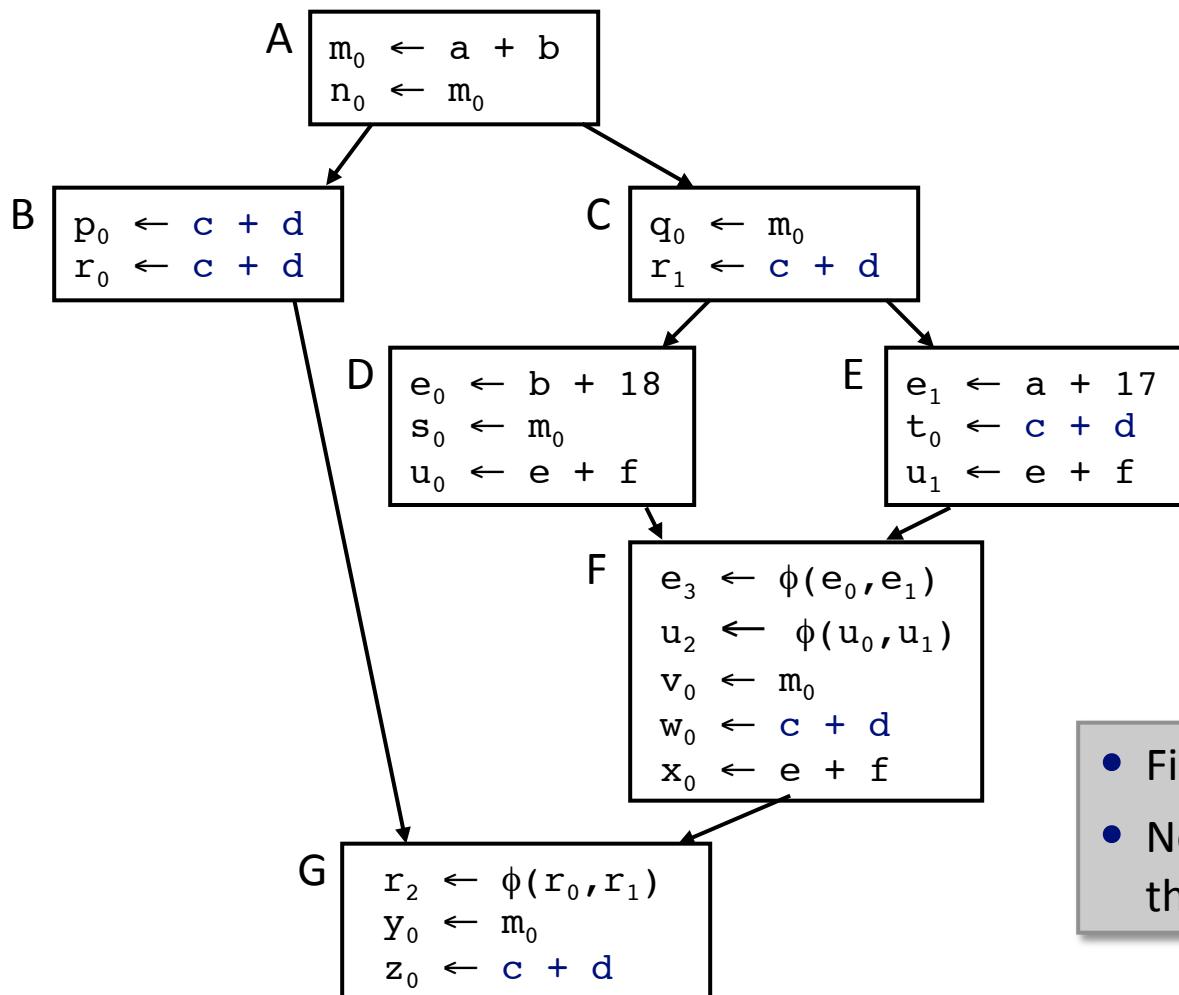
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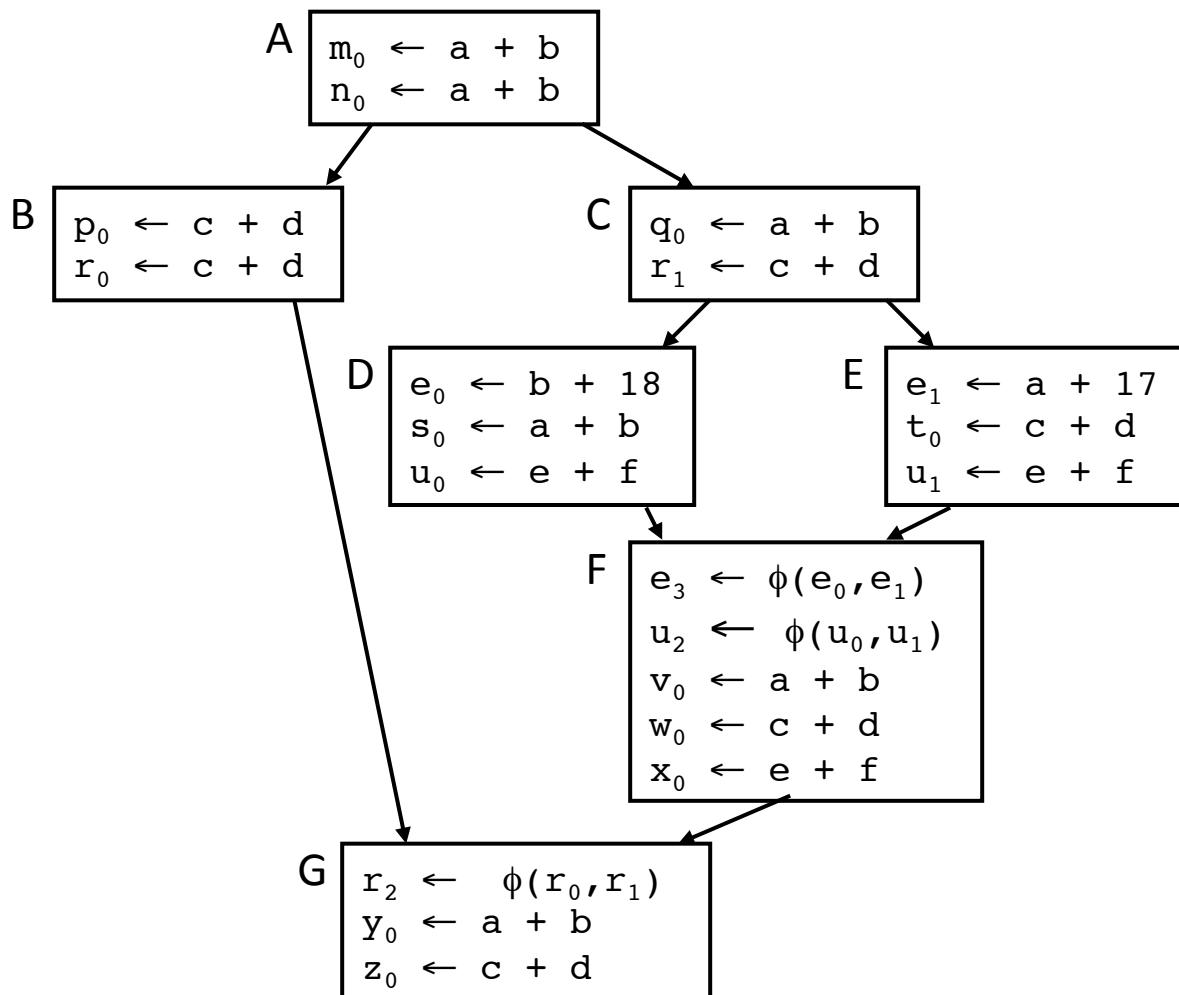
## PARTITIONING METHOD



- Fixing  $c+d$  is harder
- Need some rationale for the replacement scheme



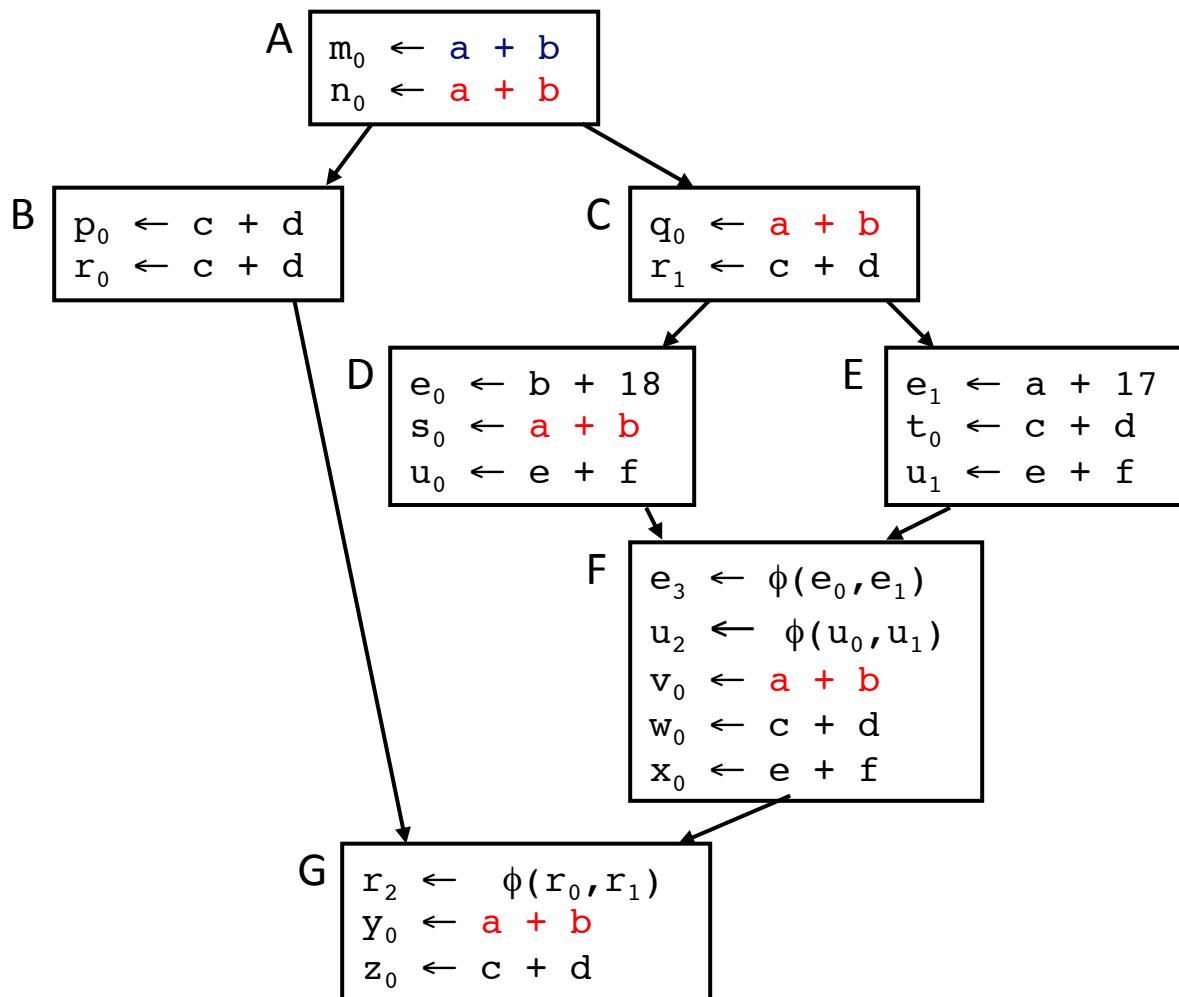
## PARTITIONING METHOD



AWZ suggest removing any computation that is dominated by another computation in the same partition



## PARTITIONING METHOD

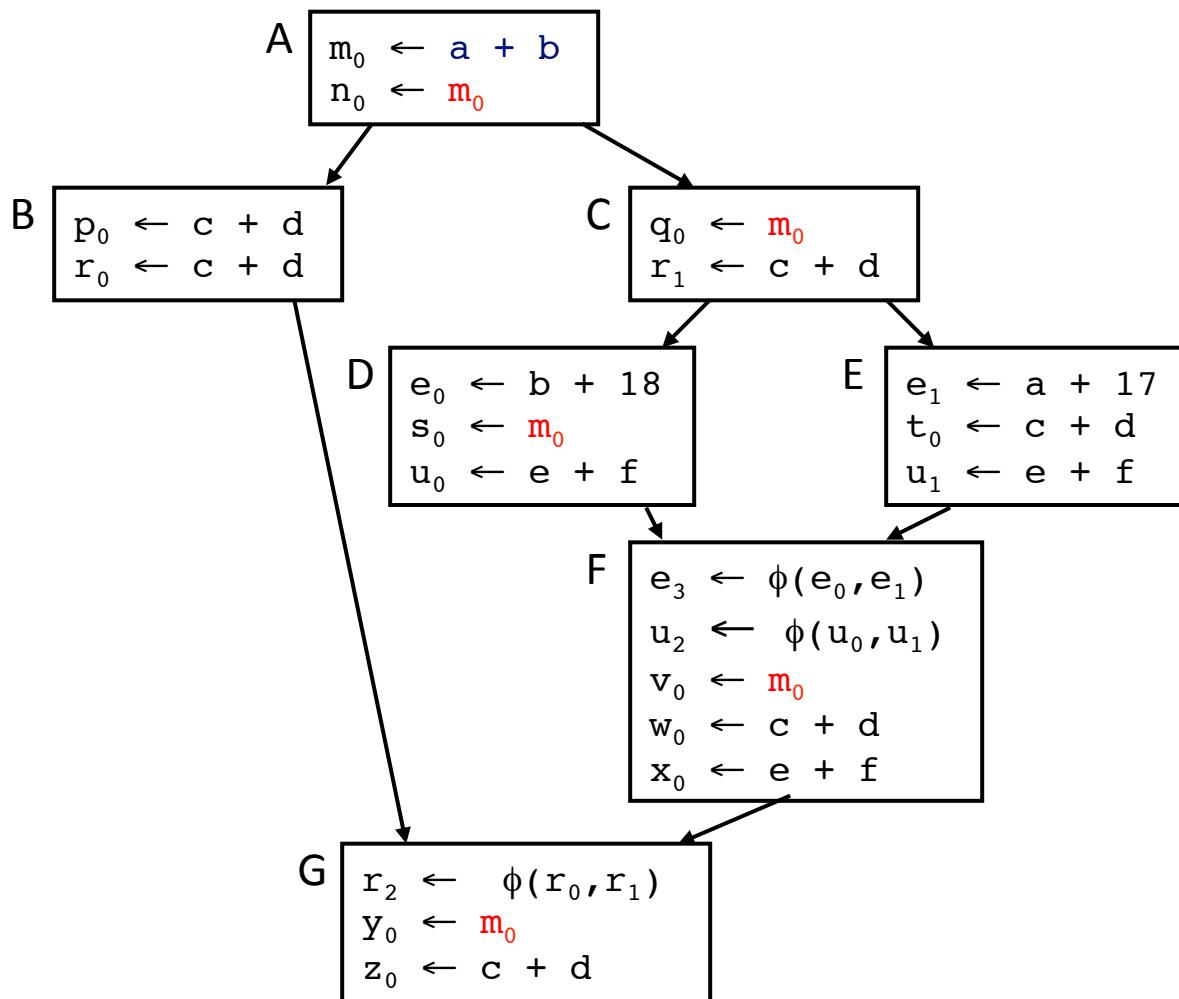


AWZ suggest removing any computation that is dominated by another computation in the same partition

- For  $a+b$ , the def of  $m_0$  dominates all other instances



## PARTITIONING METHOD

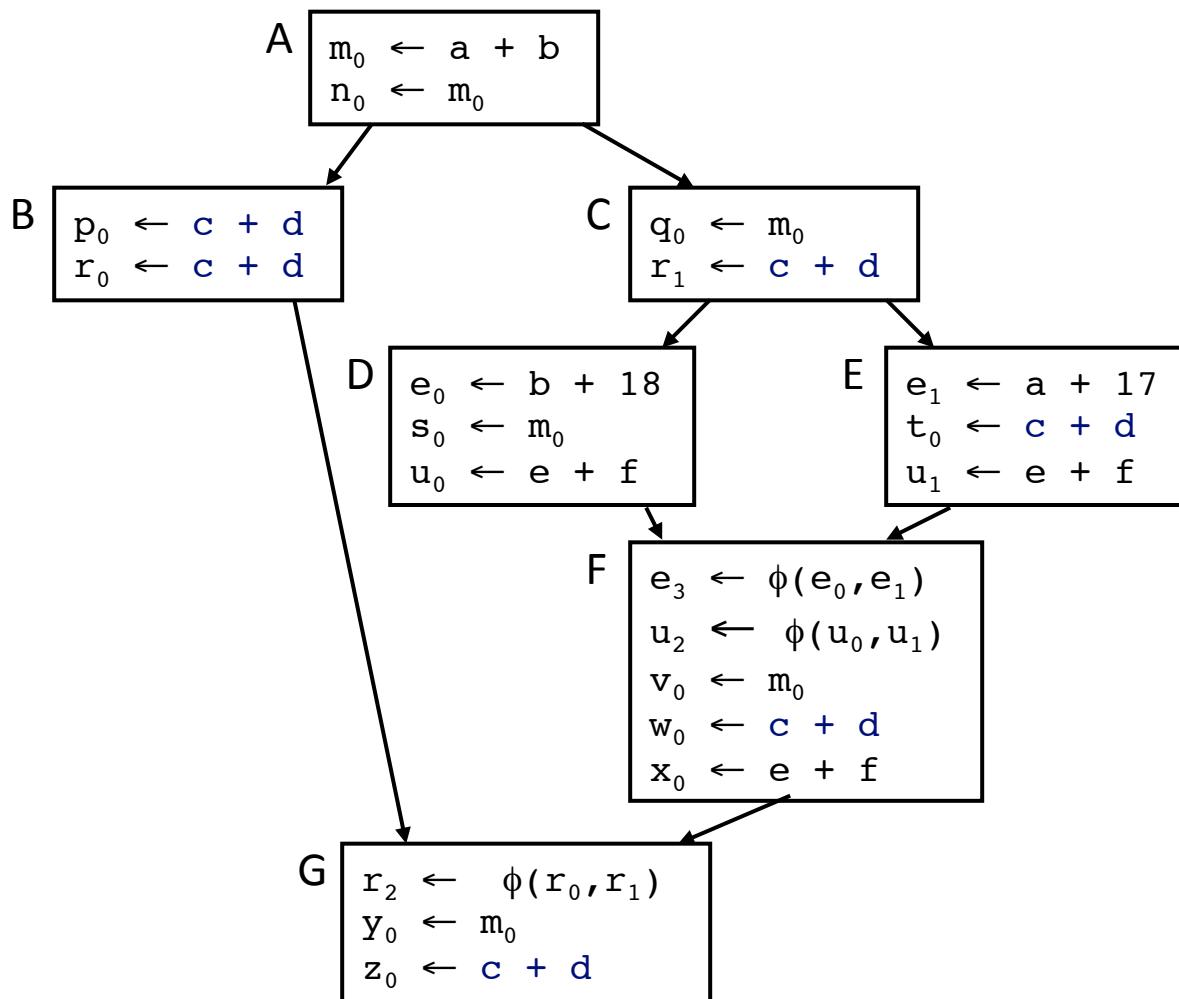


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## PARTITIONING METHOD



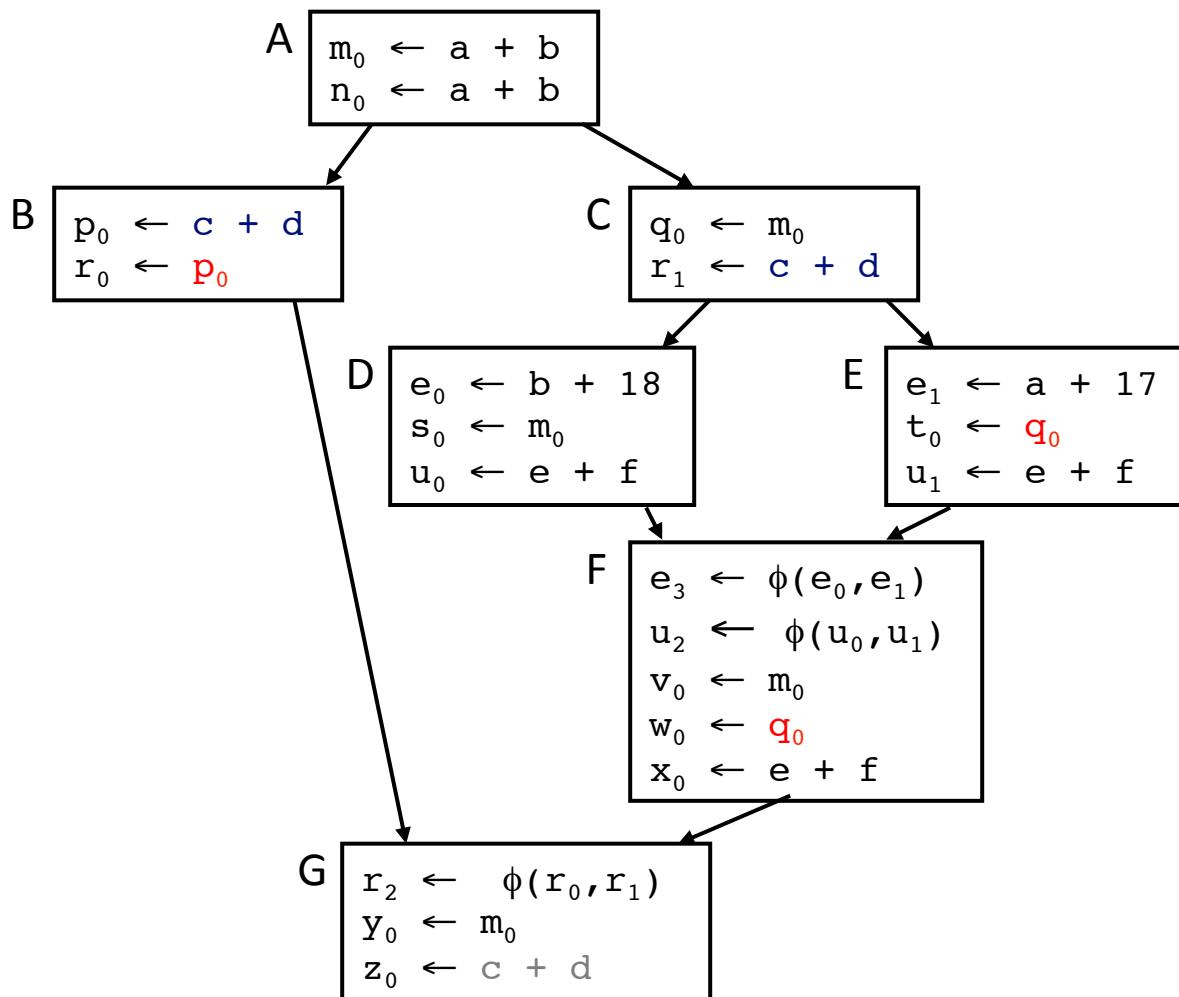
AWZ suggest removing any computation that is dominated by another computation in the same partition

- For  $a+b$ , the def of  $m_0$  dominates all other instances
- For  $c+d$  no single instance will do

\*



## PARTITIONING METHOD



AWZ suggest removing any computation that is dominated by another computation in the same partition

- For  $a+b$ , the def of  $m_0$  dominates all other instances
- For  $c+d$  no single instance will do
- Some cases cannot be caught this way

Simpson suggested another approach.



## PARTITIONING METHOD

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Simpson's suggestion

- Encode congruences into the name space
  - ◆ *This idea is critical to and derives from the Briggs-Cooper approach to algebraic reassociation for LCM/PRE*
  - ◆ Use LCM to remove redundancies

Simplification

- In SSA name space, there are no KILLS
- Knocks terms out of some equations
- Makes initial information cheaper to compute
- Larger name space, but  $O(n)$  rather than  $O(n^2)$  initialization

Consistently out-performs dominator-based replacement



## PARTITIONING METHOD

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### Is This Algorithm Optimistic Or Pessimistic?

- When we talked about optimism in data-flow analysis, we said
  - ◆ *Initialize to top*  $\Rightarrow$  optimistic
  - ◆ *Initialize to bottom*  $\Rightarrow$  pessimistic
- This algorithm does neither

AWZ builds a set of congruence relations

- Assumes maximal size sets
  - Knocks out infeasible elements
    - ◆ Over-estimates set of congruences
    - ◆ Finds largest set of self-consistent congruences
- } inherently optimistic
- ⇒ Must let it run to completion, or sets might contain lies

Click points that set construction can be either optimistic or pessimistic. An optimistic construction starts with the universe and eliminates elements. A pessimistic construction starts with the empty set and adds elements.



## PARTITIONING METHOD

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### Strengths

- Global algorithm for finding redundancies
  - ◆ *Handles arbitrary control flow by ignoring it*
  - ◆ *Finds largest congruence classes consistent with definition*
- Relies on well understood (& efficient) algorithm

### Weaknesses

- Cannot put  $2*a$  and  $a+a$  in same class (Neither can lexical methods)
- Not obvious how to handle identities or constants (Click)
- Pays no systematic attention to placement (LCM)
- Slower than DVNT (1(LVN):4(DVNT):10(AWZ))